

# Solving large linear least squares problems with equality constraints

**Jennifer Scott**

STFC Rutherford Appleton Laboratory  
and the University of Reading

**Joint with Miroslav Tůma, Charles University, Prague**

Sparse Days in St Girons 2022

# Introduction to LSE problem

Assume  $A \in \mathbb{R}^{m \times n}$  and  $C \in \mathbb{R}^{p \times n}$ , with  $m > n \gg p$ .

Further assume  $A$  is **large and sparse** and  $C$  represents a **few, possibly dense, linear constraints**.

Given  $b \in \mathbb{R}^m$  and  $d \in \mathbb{R}^p$ , the least squares problem with equality constraints (LSE problem) is

$$\min_{x \in \mathbb{R}^n} \|Ax - b\|_2 \quad \text{subject to} \quad Cx = d.$$

A solution exists if and only if  $Cx = d$  is consistent.

Solution unique if and only if  $\mathcal{A} = \begin{pmatrix} A \\ C \end{pmatrix}$  has full column rank.

## Aims of our work

Revisit solution strategies; understand their strengths and weaknesses; propose new ideas and modifications designed for tackling large-scale LSE problems.

Particular emphasis on

- ▶ the possibility that the constraints may be dense
- ▶ requiring the constraints are tightly satisfied
- ▶ testing out ideas on real data (note: published literature related to LSE problems lacks numerical results, presumably because implementing proposed methods is tough)

Only time here for a quick overview and study summary.

# Null-space approach

Classical technique (Hanson and Lawson 1969).

For LSE problems:

1. Find  $x_1 \in \mathbb{R}^n$  such that  $C x_1 = d$  (particular solution of constraint equation).
2. Construct  $Z \in \mathbb{R}^{n \times (n-p)}$  of full column rank such that  $C Z = 0$  (its columns form a basis for null-space  $\mathcal{N}(C)$ ).
3. Solve the SPD normal equations of order  $n - p$   
 $Z^T A^T A Z x_2 = (A Z)^T (b - A x_1)$ .
4. Set  $x = x_1 + Z x_2$ .

## Null-space approach: issues

- ▶ Must construct null-space basis  $Z$ .
- ▶ Traditional approach is to compute QR factorization of  $C^T$  and use to obtain  $Z$ . **But** resulting  $Z$  is dense.
- ▶ Thus  $Z^T A^T A Z$  is dense so impractical to store in memory, expensive to solve.
- ▶ **Remedy:** Scott and Tuma (2022) developed method to compute sparse (banded)  $Z$  when  $C$  is “wide” using a QR algorithm applied to  $C$  that incorporates **threshold column pivoting**.
- ▶ Threshold parameter balances stability with sparsity of  $Z$ .

## Null-space approach: results

Mixed findings: can work well but constraints not always tightly satisfied.

For these examples, constraint matrix  $C$  is dense.

$\theta$  is threshold pivoting parameter; *density* is density of  $Z^T A^T A Z$ .

Identifier	$p$	$\theta = 1$		$\theta = 0.1$	
		<i>density</i>	$\ r_c\ $	<i>density</i>	$\ r_c\ $
scagr7-2b	7	0.03	$1.27 \times 10^{-8}$	0.0007	$2.79 \times 10^{-8}$
south31	5	0.20	‡	0.02	$3.26 \times 10^{-7}$
testbig	8	0.03	$2.53 \times 10^{-11}$	0.0002	$2.90 \times 10^{-11}$
lp_fit2p	25	0.47	$4.14 \times 10^{-5}$	0.11	$3.07 \times 10^{-5}$

‡ not enough memory to run sparse direct solver on normal equs.

## Null-space approach: results

Also results can deteriorate if number of constraints increases.

Here the constraints are **not dense**.

Identifier	$p$	$\theta = 1$		$\theta = 0.1$	
		$density$	$\ r_c\ $	$density$	$\ r_c\ $
fxm4_6_5	5	0.0005	$7.80 \times 10^{-11}$	0.0005	$1.10 \times 10^{-11}$
fxm4_6_20	20	0.0006	$5.43 \times 10^{-6}$	0.0006	$6.67 \times 10^{-7}$
stormg2-8_5	5	0.002	$1.13 \times 10^{-10}$	0.002	$8.16 \times 10^{-11}$
stormg2-8_20	20	0.003	$7.44 \times 10^{-9}$	0.002	$7.23 \times 10^{-9}$

Note: if a sequence of LSE problems is to be solved with **same** set of constraints but different  $A$ , null-space basis can be reused, substantially reducing work required.

But if  $C$  changes then  $Z$  will also change (an updating strategy may be possible to reduce the work).

## Method of direct elimination (Björck and Golub 1967)

Find permutation  $P$  s.t

$$C x = C P y = \begin{pmatrix} C_1 & C_2 \end{pmatrix} \begin{pmatrix} y_1 \\ y_2 \end{pmatrix} = d,$$

with  $p \times p$  matrix  $C_1$  is **nonsingular**.

Let  $AP = \begin{pmatrix} A_1 & A_2 \end{pmatrix}$  be a conformal partitioning.

Substituting  $y_1 = C_1^{-1}(d - C_2 y_2)$  gives transformed LS problem

$$\min_{y_2} \|A_T y_2 - (b - A_1 C_1^{-1} d)\|_2^2,$$

with the **transformed** matrix

$$A_T = A_2 - A_1 C_1^{-1} C_2 \in \mathbb{R}^{m \times (n-p)}.$$



## Method of direct elimination

- ▶ If both  $A$  and  $C_1^{-1}C_2$  are sparse then solve sparse LS problem.
- ▶ If  $C$  has dense rows then  $A_{\mathcal{T}}$  has more dense rows than  $A$ .
- ▶ In this case, employ solution schemes for sparse-dense LS problems (recent work of Scott and Tůma).
- ▶ For example, Cholesky factorization of sparse part of  $A_{\mathcal{T}}$  and use as preconditioner within a CG method.
- ▶ Construct  $C_1$  using a QR factorization with a threshold parameter  $\tau$  to balance stability and limiting fill-in in  $A_{\mathcal{T}}$ .

## Method of direct elimination: results

$ndense$  is number of dense rows in  $A_T$

Identifier	$p$	$\tau = 1$		$\tau = 0.1$	
		$ndense$	$\ r_c\ $	$ndense$	$\ r_c\ $
scagr7-2b	7	14	$4.03 \times 10^{-13}$	6	$1.48 \times 10^{-13}$
scrs8-2c	22	16	$6.80 \times 10^{-14}$	16	$1.25 \times 10^{-13}$
sctap1-2b	34	72	$7.60 \times 10^{-13}$	63	$1.31 \times 10^{-12}$
gemat1_20	20	284	$8.95 \times 10^{-15}$	147	$1.64 \times 10^{-14}$
gemat1_5	5	142	$9.77 \times 10^{-14}$	17	$4.61 \times 10^{-14}$
stormg2-8_20	20	136	$5.30 \times 10^{-14}$	94	$2.29 \times 10^{-15}$
stormg2-8_5	5	61	$3.28 \times 10^{-15}$	35	$1.50 \times 10^{-14}$

- ▶  $ndense$  increases more rapidly than  $p$ .
- ▶ Reducing  $\tau$  can significantly reduce  $ndense$ , making transformed problem cheaper to solve.
- ▶ Small  $\tau$  does not adversely affect  $\|r_c\|$ .

## Method of direct elimination: weaknesses

- ▶ *ndense* can be relatively large  $\implies$  transformed problem expensive to solve.
- ▶ Suppose need to solve a sequence of problems in which  $A$  or  $C$  is fixed. Coupling of the two blocks in the solution process  $\implies$  it must be restarted for each problem.

Null-space and direct elimination approaches both face potential issues  $\implies$  other possibilities?

Consider **weighted least squares problem (WLS)**

$$\min_x \|A_\gamma x_\gamma - b_\gamma\|^2 \quad \text{with} \quad A_\gamma = \begin{pmatrix} A \\ \gamma C \end{pmatrix}, \quad b_\gamma = \begin{pmatrix} b \\ \gamma d \end{pmatrix},$$

for some large  $\gamma \gg 1$ .

$$\lim_{\gamma \rightarrow \infty} x_\gamma = x_{LSE}$$

where  $x_{LSE}$  is the solution of the LSE problem.

**Problems:** for very large values of  $\gamma$ , normal matrix  $A_\gamma^T A_\gamma$  is **extremely ill-conditioned**.

And if  $C$  has dense rows,  $A_\gamma^T A_\gamma$  is **dense**.

## Regularized normal equations

$$(A_\gamma^T A_\gamma + \omega^2 I) x = A_\gamma^T b_\gamma, \quad \omega > 0.$$

Equivalent to **3-block augmented system**

$$\begin{pmatrix} \omega I & 0 & A \\ 0 & \omega I & \gamma C \\ A^T & \gamma C^T & -\omega I \end{pmatrix} \begin{pmatrix} y_s \\ y_c \\ x \end{pmatrix} = \begin{pmatrix} b \\ \gamma d \\ 0 \end{pmatrix}$$

where  $\begin{pmatrix} y_s \\ y_c \end{pmatrix} = \omega^{-1} \left[ \begin{pmatrix} b \\ \gamma d \end{pmatrix} - \begin{pmatrix} A \\ \gamma C \end{pmatrix} x \right].$

$$\begin{pmatrix} \omega I & 0 & A \\ 0 & \omega I & \gamma C \\ A^T & \gamma C^T & -\omega I \end{pmatrix} \begin{pmatrix} y_s \\ y_c \\ x \end{pmatrix} = \begin{pmatrix} b \\ \gamma d \\ 0 \end{pmatrix}.$$

Eliminating  $y_s$  and setting  $\omega\gamma = 1$  gives  $(n+p) \times (n+p)$  system

$$\begin{pmatrix} -H(\omega) & C^T \\ C & \omega^2 I \end{pmatrix} \begin{pmatrix} x \\ y_c \end{pmatrix} = \begin{pmatrix} -A^T b \\ d \end{pmatrix}, \quad H(\omega) = A^T A + \omega^2 I.$$

Block signed (incomplete) Cholesky factorization

$$\begin{pmatrix} -H(\omega) & C^T \\ C & \omega^2 I \end{pmatrix} \approx \begin{pmatrix} L & \\ B & L_\omega \end{pmatrix} \begin{pmatrix} -I & \\ & I \end{pmatrix} \begin{pmatrix} L^T & B^T \\ & L_\omega^T \end{pmatrix},$$

$$H(\omega) \approx LL^T, \quad LB^T = -C^T, \quad S = \omega^2 I + BB^T = L_\omega L_\omega^T.$$

## Regularized normal equations: results

We use  $\omega\gamma = 1$ . GMRES convergence tolerance  $10^{-11}$

Identifier	$\omega$	$\ r\ $	Direct solver	GMRES	
			$\ r_c\ $	iters	$\ r_c\ $
sctap1-2r	$1.0 \times 10^{-3}$	$2.07 \times 10^2$	$7.38 \times 10^{-3}$	6	$7.38 \times 10^{-3}$
	$1.0 \times 10^{-5}$	$2.07 \times 10^2$	$7.71 \times 10^{-7}$	6	$7.42 \times 10^{-7}$
	$1.0 \times 10^{-9}$	$2.07 \times 10^2$	$1.29 \times 10^{-7}$	6	$4.09 \times 10^{-13}$
south31	$1.0 \times 10^{-3}$	$1.88 \times 10^2$	$8.34 \times 10^{-7}$	337	$7.34 \times 10^{-7}$
	$1.0 \times 10^{-5}$	$1.88 \times 10^2$	$8.31 \times 10^{-11}$	354	$8.85 \times 10^{-11}$
	$1.0 \times 10^{-9}$	$1.88 \times 10^2$	$6.77 \times 10^{-14}$	354	$1.08 \times 10^{-11}$
deter3_20	$1.0 \times 10^{-3}$	$1.22 \times 10^2$	$6.83 \times 10^{-6}$	36	$6.83 \times 10^{-6}$
	$1.0 \times 10^{-5}$	$1.22 \times 10^2$	$6.83 \times 10^{-10}$	36	$6.83 \times 10^{-10}$
	$1.0 \times 10^{-9}$	$1.22 \times 10^2$	$1.35 \times 10^{-12}$	36	$1.37 \times 10^{-12}$

## Comments on using augmented system

- ▶ Several options for using an augmented system formulation (just one considered here).
- ▶ Allows use of “black box” sparse solvers  $\implies$  greatly reduces work to develop robust and efficient implementations.
- ▶ Can be generalised to handle dense rows in  $A$  and offers the potential for mixed-precision computation.
- ▶ Significant amount of work can be reused when solving sequence of problems in which only  $C$  changes.



## Concluding remarks

- ▶ Large-scale LSEs can be tough to solve, particularly if constraints are dense.
- ▶ We need to bring together different techniques and experience with LS problems (especially for mixed sparse-dense case).
- ▶ General lack of iterative methods and preconditioners that can be used to extend the size of problems that can be solved.
- ▶ Software packages?

# Thank you for listening!

Scott and Tũma (2022) *Solving large linear least squares problems with linear equality constraints*.

To appear in BIT. Also <https://arxiv.org/abs/2106.13142>

Papers on relating to our work on sparse LS problems available from [http://www.numerical.rl.ac.uk/people/j\\_scott/](http://www.numerical.rl.ac.uk/people/j_scott/)

My work was partially supported by EPSRC grant EP/I013067/1

## Questions or comments?